

An Analysis of algebraic preconditioning based on different variants of the Gram-Schmidt algorithm

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1. 10. 2010

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Motivation of preconditioning

$$\mathbf{Ax} = \mathbf{b}$$

- improving of spectral properties of the matrix \mathbf{A}
- absolute necessity for increasing robustness of iterative solvers
- acceleration of convergence

General strategy: multiply equation $\mathbf{Ax} = \mathbf{b}$ by \mathbf{P} , where $\mathbf{P} \approx \mathbf{A}^{-1}$, so that $\|\mathbf{I} - \mathbf{PA}\| \rightarrow 0$

Preserving symmetry

Assume A is SPD

For preserving symmetry: $\mathbf{P} = \mathbf{Z}\mathbf{Z}^T$

$$\|\mathbf{I} - \mathbf{Z}^T \mathbf{A} \mathbf{Z}\| \rightarrow 0$$

Question: How to compute preconditioning \mathbf{Z} efficiently, so that computation cost is much less $O(n^3)$ (Gaussian elimination) for dense case and/or parallelizable?

Algorithms for preconditioning

- ILU - classical type of preconditioning
- SPAI - minimization of functional $\|\mathbf{I} - \mathbf{PA}\|_F$, decomposition to n-independent problems, TU Munich
- AINV, SAINV - based on the generalized Gram-Schmidt algorithm
- ...

Preconditioning matrix \mathbf{Z} should be sparse, in order to have $\mathbf{Z}^T \mathbf{A} \mathbf{Z}$ sparse as well \rightarrow prescribing pattern, dropping...

Successful preconditioning

Accuracy: $\|\mathbf{I} - \mathbf{Z}^T \mathbf{A} \mathbf{Z}\|$

Stability: $\|\mathbf{A} - (\mathbf{Z}^T \mathbf{Z})^{-1}\|$

- Low
- Controlled (via dropping)

NOTE: Currently no general theory of preconditioning
(stability + accuracy) \rightarrow convergence

Generalized Gram-Schmidt algorithm

- \mathbf{A} is a SPD matrix of dimension $n \times n$
- energetic dot product $\langle x, y \rangle_A = y^T \mathbf{A} x$
- \mathbf{A} -orthogonality (conjugate gradient method)
- basis $\mathbf{Z}^{(0)}$, which will be \mathbf{A} -orthogonalized against previously computed vectors
(for simplicity and for real-world problems $\mathbf{Z}^{(0)} = \mathbf{I}$ or $\mathbf{Z}^{(0)} = \text{diag}(\mathbf{A})^{-1/2}$)

Versions

- classical, modified, with iterative refinement
- right looking, left looking

Properties in exact arithmetic

- $\mathbf{Z}^T \mathbf{A} \mathbf{Z} = \mathbf{I}$
- $\mathbf{A}^{-1} = \mathbf{Z} \mathbf{Z}^T$
- $\mathbf{Z}^{(0)} = \mathbf{Z} \mathbf{U}$
- $(\mathbf{Z}^{(0)})^T \mathbf{A} \mathbf{Z}^{(0)} = \mathbf{U}^T \mathbf{U}$

Finite precision

- In finite precision arithmetic $\|\mathbf{Z}^T \mathbf{A} \mathbf{Z} - \mathbf{I}\| < ???$

Question: How to estimate the loss of orthogonality $\|\mathbf{Z}^T \mathbf{A} \mathbf{Z} - \mathbf{I}\|$ in case of incomplete Gram-Schmidt algorithm (with dropping)?

Idea: Björck, Å., *Solving linear least squares problems by Gram-Schmidt orthogonalization*, BIT 7 (1967), 1-21

Preconditioning theory

Analysis:

Behavior in finite precision arithmetic (rounding errors)

Practical algorithms:

Dropping and incomplete algorithm Gram-Schmidt

Theory: Move from (full) algorithm behavior (u) to preconditioning (incomplete algorithm - τ)

Our results

$$\begin{aligned}\|\mathbf{Z}^T \mathbf{A} \mathbf{Z} - \mathbf{I}\| &\leq O(u) \kappa(A) && (\text{CGS2, EIG}) \\ &\leq \frac{O(u) \kappa^{3/2}(A)}{1 - O(u) \kappa^{3/2}(A)} && (\text{MGS - SAINV}) \\ &\leq \frac{O(u) \kappa^2(A)}{1 - O(u) \kappa(A)} && (\text{CGS, AINV})\end{aligned}$$

Choice of $\mathbf{Z}^{(0)}$

$\mathbf{Z}^{(0)}$ can be any $n \times m$ matrix with full column rank and $m \leq n$
 (\mathbf{A} is an $n \times n$ matrix)

Loss of orthogonality (MGS):

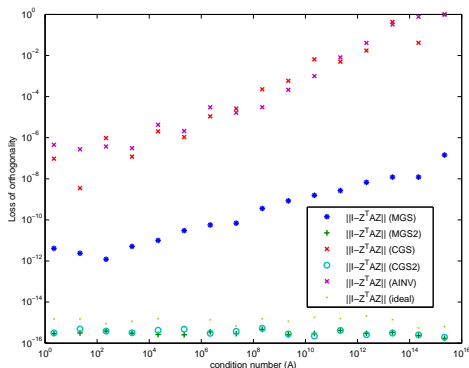
$$\|\mathbf{Z}^T \mathbf{A} \mathbf{Z} - \mathbf{I}\| \leq \frac{O(u) \kappa(\mathbf{A}) \kappa(\mathbf{A}^{1/2} \mathbf{Z}_k^{(0)})}{1 - O(u) \kappa(\mathbf{A}) \kappa(\mathbf{A}^{1/2} \mathbf{Z}_k^{(0)})}$$

NOTE: If $\mathbf{Z}^{(0)}$ is an upper triangular matrix, then $\mathbf{Z} = \mathbf{L}^{-1}$, where $\mathbf{A} = \mathbf{L} \mathbf{L}^T$ (Cholesky factorization), otherwise \mathbf{Z} is general matrix.

Test problem 1

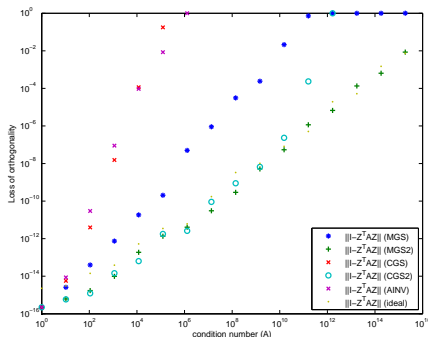
\mathbf{A} = diagonal matrix, $\kappa(\mathbf{A}) \approx 10^i, i = 0, \dots, 15$

$\mathbf{Z}^{(0)} = \mathbf{H}^{1/2}$, where \mathbf{H} is the inverse Hilbert matrix, $\kappa(\mathbf{Z}^{(0)}) \approx 10^5$



Test problem 2

$\mathbf{A}, \mathbf{Z}^{(0)}$ are the inverse Hilbert matrices (general powers)
 $\kappa(\mathbf{A}), \kappa(\mathbf{Z}^{(0)}) \approx 10^i, i = 0, \dots, 15$



Conclusion and open questions

Done

- (full) algorithm analysis

Future work

- find a connection between estimate of $\|\mathbf{Z}^T \mathbf{A} \mathbf{Z} - \mathbf{I}\|$ and dropping strategy
- analyze behavior of preconditioned iterative methods
- analyze other related algorithms and preconditioners

Thank you for your attention!!!